# Combinatorial proof of the Graham-Pollak Determinant formula for the distance matrix of a tree* 

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#### Abstract

Graham and Pollak proved in 1971 that the determinant of the distance matrix of any tree with $n$ vertices is $(-1)^{m-1}(n-1) 2^{n-2}$. We provide a combinatorial interpretation for this formula, and, at the same time, the first combinatorial proof for this formula. Our approach is based on the Gessel-Viennot-Lindström Lemma. Our methods suits as well for many parametric deformations of the Graham-Pollak Formula. The special case when the tree is a path graph allows to recover results on statistics on derangements.


## 1 Introduction

Graham and Pollak [7] proved in 1971 the following surprising theorem: the determinant of the distance matrix of any tree with $n$ vertices is independent on its structure, and always equal to

$$
(-1)^{n-1}(n-1) 2^{n-2}
$$

This result is not difficult to prove using row and column operations, and induction. Yet, such a formula strongly suggests that this determinant counts something. But what? We answer this question (raised for instance in $[9,12]$ ), and also provide a combinatorial proof of the formula.

Evaluating the Graham-Pollak determinant is easily seen to be a signed enumeration problem, i.e. it consists in simplifying a sum:

$$
\sum_{x \in X} \varepsilon(x)
$$

where $x$ are the elements of some set $X$ and $\varepsilon$ is a sign function (taking values +1 and -1 ). Here the set $X$ is the set of all pairs $(\sigma, f)$ where $\sigma$ is a permutation of the vertices of the tree; and $f$ is a function assigning to each vertex $v$ an edge $f(v)$ in the unique path from $v$ to $\sigma(v)$. We call these pairs the compatible pairs for the tree $T$.

A standard strategy to deal with signed enumeration problems consists in exhibiting a signreversing involution on $X$, i.e. an involution $\tau$ of $X$ such that (i) if $x$ is not a fixed point of $\tau$ then

[^0]its image has opposite sign and (ii) all fixed points $x$ have the same sign. Then most terms in the signed sum cancel, and the sum simplifies into the cardinality of the set fixed points, up to the sign.

A further refinement of the sign-reversing involution strategy, based on the Gessel-ViennotLindström Lemma [6] (see also the exposition [1, Ch. 32 (Lattice paths and determinants)]), consists in interpreting each element of $X$ as a $n$-path (a family of $n$-paths from $n$ fixed sources to $n$ fixed sinks) in a acyclic digraph $G$. Any such $n$-path matches the sources with the sinks and defines therefore a permutation of $\{1,2 \ldots, n\}$, whose sign is the sign of the $n$-path. Then the signed sum of these $n$-paths is the number of intersection free $n$-paths. Indeed, there is a signed reversing involution that kills the $n$-paths that have an intersection.

## 2 Main results and sketch of the proofs

Let $T$ be a tree with $n$ vertices. Let $X$ be the set of its compatible pairs $(\sigma, f)$.

1. Define a diagram on $T$ as a multiset $D$ of $n$ oriented edges of the tree $T$. To each compatible pair $(\sigma, f)$ we associate the following diagram $D$ : the multiset of the edges $f(v)$, where each $f(v)$ is oriented in the same direction as the path from $v$ to $\sigma(v)$. This splits $X$ into a disjoint union of classes $X(D)$, and therefore

$$
\sum_{x \in X} \varepsilon(x)=\sum_{D} \sum_{x \in X(D)} \varepsilon(x)
$$

2. Say that a diagram is standard if for each edge $x-y$ of the tree $T$, there is one oriented edge $(x \rightarrow y$ or $x \leftarrow y)$ in $D$, except for one, for which both oriented edges $(x \rightarrow y$ and $x \leftarrow y)$ appear. Note that there are $(n-1) 2^{n-2}$ standard diagrams. Indeed, the tree $T$ has $n-1$ edges; there are $n-1$ choices for the edge with double orientation, and $2^{n-2}$ choices of orientations for the remaining $n-2$ edges.
3. For $D$ non-standard, we exhibit a sign-reversing involution on $X(D)$ with no fixed point. Its existence shows that:

$$
\sum_{x \in X(D)} \varepsilon(x)=0
$$

4. For $D$ standard, we build an acyclic digraph $G(D)$ with $n$ sources and $n$ sinks, such that the elements of $X(D)$ are in a sign-preserving bijection with the $n$-paths in $G(D)$. We show that there is only one intersection-free $n-$ path in $G(D)$, and the corresponding permutation is a full cycle. Therefore,

$$
\sum_{x \in X(D)} \varepsilon(x)=(-1)^{n-1}
$$

5. This finishes the combinatorial proof and the combinatorial interpretation of the formula: $(n-1) 2^{n-2}$ counts the standard diagrams and $(-1)^{n-1}$ is the sign of a full cyclic permutation of $\{1,2, \ldots, n\}$.

## 3 Parametric deformations

Several deformations of the Graham-Pollak Formula, obtained by introducing parameters, have been obtained $[3-5,10,11]$.

For instance, when the distances are replaced with their $q$-analogues (each distance $d$ is replaced with the corresponding " $q$-integer" $1+q+\cdots+q^{d-1}$, where $q$ is a variable), the determinant of the distance matrix of any tree becomes, as shown in [10, corollary 2.3],

$$
\begin{equation*}
(-1)^{n-1}(n-1)(q+1)^{n-2} \tag{1}
\end{equation*}
$$

Another example comes from attaching to each edge $e$ a variable $x_{e}$. In the determinant of the distance matrix, replace the distance between vertices $v$ and $w$ with $\sum x_{e}$, where the sum is over all edges in the path from $v$ to $w$. Then the determinant of the distance matrix becomes, as shown in [3, Corollary 2.5],

$$
\begin{equation*}
(-1)^{n-1} 2^{n-2} \sum_{e} x_{e} \prod_{e} x_{e} \tag{2}
\end{equation*}
$$

Our approach fits like a glove for the study of these deformations. Indeed, the Gessel-ViennotLindström Lemma applies to weighted acyclic digraphs, and serves to count weighted signed sums. We recover many of the deformations of the Graham-Pollak Formula found in the literature (and all of $[3-5,10,11]$ ), simply by equipping the graphs $G(D)$ with suitable weights.

## 4 Derangements

We apply the same strategy to problems of signed enumeration of derangements (permutations with no fixed point) . In [8] it is proved that:

$$
\begin{equation*}
\sum_{\sigma} \varepsilon(\sigma) q^{\operatorname{exc}(\sigma)}=(-1)^{n-1}\left(q+q^{2}+\cdots+q^{n-1}\right) \tag{3}
\end{equation*}
$$

where the sum is over all derangements $\sigma$ of $\{1,2, \ldots, n\}, \operatorname{exc}(\sigma)$ is the number of excedances of $\sigma$ (the number of indices $i$ such that $\sigma(i)>i$ ) and $\varepsilon(\sigma)$ is the sign of the permutation $\sigma$.

The identity (3) was refined recently in [2, Formula (3)] as:

$$
\begin{equation*}
\sum_{\sigma} \varepsilon(\sigma)\left(\prod_{j \in \operatorname{RLM}(\sigma)} x_{j}\right)\left(\prod_{j \in \operatorname{EXC}(\sigma)} y_{j}\right)=(-1)^{n-1} \sum_{j=1}^{n-1} y_{1} \cdots y_{j} x_{j+1} \cdots x_{n} \tag{4}
\end{equation*}
$$

where $\operatorname{RLM}(\sigma)$ is the set of right-to-left minima of $\sigma$ (the set of all indices $i$ such that $\sigma(i)<\sigma(j)$ for all $j>i$ ) and $\operatorname{EXC}(\sigma)$ is the set of excedances of $\sigma$ (the set of all indices $i$ such that $\sigma(i)>i)$.

Again, we turn the sign enumeration problems (raised by the left-hand side of the formulas) into problems of non-intersecting paths in weighted acyclic digraphs with $n$ sources and $n$ sinks. We find ( $n-1$ ) non-intersecting $n$-paths, that correspond to $n-1$ terms of the right-hand sides of (3) and (4). The proof is remarkably simple and visually appealing.

## 5 Further remarks

We observe that many proofs with sign-reversing involutions involve arbitrary choices (in the definition of the involution). Indeed, often the sign-reversing involution is not unique. This sometimes makes the description of the involution, and the subsequent proofs, cumbersome. Turning the problem into a non-intersecting $n$-path problem avoids this difficulty, by changing the problem of constructing an involution (one among many), into a problem of finding a set of non-intersecting $n$-paths (a well-defined set). This provides clearer proofs.

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